Introduction to Discrete-time Averaging Systems

Chengcheng Zhao

Zhejiang University, Hangzhou, China &University of Victoria, Victoria, Canada

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Discrete-time Averaging Systems

August 2, 2021 1 / 32

• Swarm behaviour in nature



(a) Bird flocking



(b) Ant swarming



(c) Fish Swarming

• Swarm behaviour in nature



(d) Bird flocking

(e) Ant swarming

(f) Fish Swarming

- No control center
- Individual animals only interact with their neighbours
- Collective animal behavior

How can we use the idea behind in social and engineering fields?

Consensus algorithms

- Average consensus: all states converges to average
- Maximum consensus: all states converge to maximum value
- Wide application
 - Smart grids, VANETS, social networks, and crowd-sensing



Figure 1: Wide applications

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Figure 1: Wide applications

Two application examples to show the related averaging systems

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August 2, 2021 3 / 32





- Social influence networks: opinion dynamics
 - $\bullet\,$ A group of n individuals who must act together as a team



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 - Each individual has its own subjective estimate $p_i \mbox{ for the unknown parameters }$
 - Individual i is appraised of p_i of each other member $j \neq i$ of the group
 - How to model predictions that the individual will revise its estimate?

- Social influence networks: opinion dynamics
 - The model (French-Harary-DeGroot) predicts that the individual will revise its estimate to be

$$p_i(k+1) = \sum_{j=1}^n a_{ij} p_j(k)$$
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$$p_i(k+1) = \sum_{j=1}^n a_{ij} p_j(k)$$
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- $a_{ij} \ge 0$ denotes the weight that individual *i* assigns to individual *j*;
- $\sum_{j=1}^{} a_{ij} = 1$ for all i;
- a_{ii} describes the attachment of individual i to its own opinion;
- a_{ij} is an interpersonal influence weight that individual i accords to individual j;

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- Scientific questions of interests
 - Is this model of human opinion dynamics believable? Is there empirical evidence in its support?
 - How does one measure the coefficients a_{ij} ?
 - Under what conditions do the estimate converge to the same estimate? And to what final estimate?
 - What are more realistic, empirically-motivated models, possibly including stubborn individuals or antagonistic interactions?



- Wireless sensor networks
 - A collection of spatially-distributed devices



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How can all devices obtain the accurate estimate in a distributed way?



Figure 3: The communication graph for devices

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 - Each node has a measured temperature $x_i(0)$



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$$x_1(k+1) = x_1(k)/2 + x_2(k)/2$$

• update rule x(k+1) = Ax(k)

$$\begin{bmatrix} x_1(k+1)\\ x_2(k+1)\\ x_3(k+1)\\ x_4(k+1) \end{bmatrix} = \begin{bmatrix} 1/2 & 1/2 & 0 & 0\\ 1/4 & 1/4 & 1/4 & 1/4\\ 0 & 1/3 & 1/3 & 1/3\\ 0 & 1/3 & 1/3 & 1/3 \end{bmatrix} \begin{bmatrix} x_1(k)\\ x_2(k)\\ x_3(k)\\ x_4(k) \end{bmatrix}$$

(3)

• Apply Algorithm 1

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Image: A math a math

(4)

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Figure 4: The original communication graph and the weighted graph

(4)

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(5)

Apply Algorithm 1



As the average is 24, average consensus cannot be achieved.

(5)

• Apply a new weight strategy (Algorithm 2)

$$\begin{bmatrix} x_1(k+1)\\ x_2(k+1)\\ x_3(k+1)\\ x_4(k+1) \end{bmatrix} = \begin{bmatrix} 3/4 & 1/4 & 0 & 0\\ 1/4 & 1/4 & 1/4 & 1/4\\ 0 & 1/4 & 5/12 & 1/3\\ 0 & 1/4 & 1/3 & 5/12 \end{bmatrix} \begin{bmatrix} x_1(k)\\ x_2(k)\\ x_3(k)\\ x_4(k) \end{bmatrix}$$

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Figure 5: The original communication graph and the weighted graph

11 / 32

(6)

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As the average is 24, average consensus can be achieved.

(7)

$$A = \begin{bmatrix} 3/4 & 1/4 & 0 & 0\\ 1/4 & 1/4 & 1/4 & 1/4\\ 0 & 1/4 & 5/12 & 1/3\\ 0 & 1/4 & 1/3 & 5/12 \end{bmatrix}, A = \begin{bmatrix} 1/2 & 1/2 & 0 & 0\\ 1/4 & 1/4 & 1/4 & 1/4\\ 0 & 1/3 & 1/3 & 1/3\\ 0 & 1/3 & 1/3 & 1/3 \end{bmatrix}$$

• A is a non-negative matrix

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- A is a non-negative matrix
- A is a row stochastic matrix

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- $\bullet\,$ The associated graph of A is strongly connected

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- When can we achieve average consensus?

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- A is a non-negative matrix
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- When can we achieve average consensus? In Algorithm 2, A is symmetric

• Dynamic model

$$x(k+1) = Ax(k) \tag{9}$$

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$$x(k+1) = Ax(k)$$

$$a_{12} \cdots a_{1n}$$

$$(9)$$

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$$A = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{1n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{nn} \end{bmatrix}, x(k) = \begin{bmatrix} x_1(k) \\ x_2(k) \\ \vdots \\ x_n(k) \end{bmatrix}$$
(10)

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- $A \in \mathbb{R}^{n \times n}$ has non-negative entries and unit row sums
- $x(k) \in \mathbb{R}^n$, $k \ge 0$ $x_i(0)$ is the initial scalar state (temperature, vibrations, sound, light) $x_i(k)$ is the updated state at iteration k

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- Interesting problems for the averaging model
 - Does each node converge to a value? Is this value the same for all nodes (consensus)?

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• Interesting problems for the averaging model

- Does each node converge to a value? Is this value the same for all nodes (consensus)?
- Is this value equal to the average of the initial conditions? When do the agents achieve average consensus?
- What properties do the graph and the corresponding matrix need to have in order for the algorithm to converge?

• Dynamic model $x(k+1) = Ax(k) \Rightarrow x(k) = Ax(k-1)$ $= A \times Ax(k-1)$ $= \vdots$ $= A^{(k+1)}x(0)$ (11)

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Jordan normal form

$$A = PJP^{-1} \Rightarrow x(k+1) = A^{(k+1)}x(0)$$

= $(PJP^{-1})^{k+1}x(0)$ (12)

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$$A^{2} = PJP^{-1}PJP^{-1} = PJ^{2}P^{-1}$$

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(14)

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(14)

• Transformation

$$J = \begin{bmatrix} \lambda_1 & \cdots & 0\\ \vdots & \ddots & \vdots\\ 0 & \cdots & \lambda_n \end{bmatrix}$$
$$\Rightarrow J^{k+1} = \begin{bmatrix} \lambda_1^{k+1} & \cdots & 0\\ \vdots & \ddots & \vdots\\ 0 & \cdots & \lambda_n^{k+1} \end{bmatrix}$$

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• Suppose $A \in \mathbb{R}^{n \times n}$ is symmetric

• Take limitations on both sides of the equation

$$\lim_{k \to \infty} x(k+1) = \lim_{k \to \infty} P J^{k+1} P^{-1} x(0)$$

$$= \lim_{k \to \infty} P \begin{bmatrix} \lambda_1^{k+1} & \cdots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \cdots & \lambda_n^{k+1} \end{bmatrix} P^{-1} x(0)$$

$$= P \begin{bmatrix} \lim_{k \to \infty} \lambda_1^{k+1} & \cdots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \cdots & \lim_{k \to \infty} \lambda_n^{k+1} \end{bmatrix} P^{-1} x(0)$$
(15)

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$$= P \begin{bmatrix} \lim_{k \to \infty} \lambda_1^{k+1} & \cdots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \cdots & \lim_{k \to \infty} \lambda_n^{k+1} \end{bmatrix} P^{-1} x(0)$$
(15)

• Consensus is correlated to the eigenvalues of the matrix A • Limitation exists if $\lim_{k\to\infty}\lambda_i^{k+1}$ exists, i.e., $\lambda_i\leq 1$

• The power of matrix A

$$A^{2} = \begin{bmatrix} 1/2 & 1/2 & 0 & 0 \\ 1/4 & 1/4 & 1/4 & 1/4 \\ 0 & 1/3 & 1/3 & 1/3 \\ 0 & 1/3 & 1/3 & 1/3 \end{bmatrix} \begin{bmatrix} 1/2 & 1/2 & 0 & 0 \\ 1/4 & 1/4 & 1/4 & 1/4 \\ 0 & 1/3 & 1/3 & 1/3 \end{bmatrix} = \begin{bmatrix} 0.3750 & 0.3750 & 0.1250 & 0.1250 \\ 0.1875 & 0.3542 & 0.2292 & 0.2292 \\ 0.0833 & 0.3056 & 0.3056 & 0.30560 \\ 0.0833 & 0.3056 & 0.3056 & 0.3056 \end{bmatrix}$$
(16)

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(16)



Figure 6: The original communication graph and the weighted graph

- The power of matrix A
 - $\bullet\,$ Nonzero elements of $A^2:$ the directed path with a length of 2 in the associated graph
 - $[A^2]_{ij} > 0$, there is a directed path between node i and node j
 - The information flow between different nodes $[A^2]_{ij} > 0$, node i can obtain the information of node j through two hops interaction

Row-stochastic matrices and their spectral radius

- For any row-stochastic matrix $A \in \mathbb{R}^{n \times n}$
 - 1) 1 is an eigenvalue \Leftarrow definition $A1_n = 1_n$
 - 2) spec(A) is a subset of the unit disk and $\rho(A) = 1$
- Gershgorin Disk Theorem

Theorem

For any square matrix $A \in \mathbb{R}^{n \times n}$,

$$\operatorname{spec}(A) \subset \bigcup_{i=\{1,\dots,n\}} \{ z || z - a_{ii} | \le \sum_{j=1, j \neq i}^{n} |a_{ij}| \}$$
 (17)

Proof.

$$Ax = \lambda x, \ x \neq 0_n, \ |x_i| = \max_{j\{1, \dots, n\}} |x_j| > 0 \Rightarrow \lambda x_i = \sum_{j=1}^n a_{ij} x_j$$
$$\Rightarrow \lambda - a_{ii} = \sum_{j=1, j \neq i}^n a_{ij} x_j / x_i$$
$$\Rightarrow |\lambda - a_{ii}| = |\sum_{j=1, j \neq i}^n a_{ij} x_j / x_i| \le \sum_{j=1, j \neq i}^n |a_{ij}| |x_j| / |x_i| \le \sum_{j=1, j \neq i}^n |a_{ij}| \qquad \square$$

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Discrete-time Averaging Systems

Perron-Frobenius Theory

- Irreducible and primitive matrices
 - $A \in \mathbb{R}^{n \times n}$, $n \geq 2$ has non-negative entries and is
 - $\bullet~~{\rm irreducible}~{\rm if}~\sum\limits_{k=0}^{n-1}A^k>0$ (G is strongly connected)
 - primitive if there exists a positive integer k such that $A^k > 0$ (G is strongly connected and aperiodic)
 - a primitive matrix is irreducible



Figure 7: The set of non-negative square matrices and its subsets of irreducible, primitive and positive matrices

Irreducible and primitive matrices

$$\begin{array}{ll} A_1 = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} & : \operatorname{spec}(A_1) = \{1, 1\}, \mbox{ the zero/nonzero pattern in } A_1^k \mbox{ is constant, and } \\ \lim_{k \to \infty} A_1^k = I_2, \\ A_2 = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} & : \operatorname{spec}(A_2) = \{1, -1\}, \mbox{ the zero/nonzero pattern in } A_2^k \mbox{ is periodic, and } \\ \lim_{k \to \infty} A_2^k \mbox{ does not exist,} \\ A_3 = \begin{bmatrix} 0 & 1 \\ 0 & 0 \end{bmatrix} & : \operatorname{spec}(A_3) = \{0, 0\}, \mbox{ the zero/nonzero pattern is } A_3^k = 0 \mbox{ for all } k \ge 2, \mbox{ and } \\ \lim_{k \to \infty} A_3^k = 0, \\ A_4 = \begin{bmatrix} \frac{1}{2} & \frac{1}{2} \\ 1 & 0 \end{bmatrix} & : \operatorname{spec}(A_4) = \{1, -1/2\}, \mbox{ the zero/nonzero pattern is } A_4^k > 0 \mbox{ for all } k \ge 2, \mbox{ and } \\ \lim_{k \to \infty} A_4^k = 0, \\ A_5 = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} & : \operatorname{spec}(A_5) = \{1, 1\}, \mbox{ the zero/nonzero pattern in } A_5^k \mbox{ is constant and } \\ \lim_{k \to \infty} A_5^k \mbox{ is unbounded.} \end{array}$$

Figure 8: Example 2-dimensional non-negative matrices and their properties

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Perron-Frobenius Theorem

Theorem

Let $A \in \mathbb{R}^{n \times n}$, $n \ge 2$. If A is non-negative, then 1) there exists a real eigenvalue $\lambda \ge |\mu| \ge 0$ for all other eigenvalues μ ; 2) the right and left eigenvectors v and w of λ can be selected non-negative. If additionally A is irreducible, then 3) the eigenvalue λ is strictly positive and simple; 4) the right and left eigenvectors v and w of λ are unique and positive. If additionally A is primitive, then 5) the eigenvalue $\lambda \ge |\mu|$ for all other eigenvalues μ

• Proof: analyze properties of positive matrices and then use "limit"

Perron-Frobenius Theory

• Lemma for positive matrices

Lemma

Let $A \in \mathbb{R}^{n \times n}$, $n \ge 2$. If A is positive, then Lem-1) there exists an eigenvalue $\lambda = \rho(A) > |\mu| \ge 0$ for all other eigenvalues μ ; Lem-2) λ is simple, i.e., $\operatorname{algmulti}_A(\lambda) = 1$; Lem-3) the right and left eigenvectors v and w of λ are positive.

- Proof is omitted and can be found in the reference below
- $\rho(A)$ is the only one eigenvalue on the spectral circle
- $\bullet\,$ Algebraic and geometric multiples are equal to 1

C. D. Meyer, "Matrix analysis and applied linear algebra," SIAM, 2000.

Perron-Frobenius Theory

- Proof of 1) and 2) of non-negative matrix A
 - Key idea: positive matrices \Rightarrow sequence convergence
 - Construct a positive matrix $A_k = A + (1/k) \mathbf{1}_n \mathbf{1}_n^\top$
 - $\Rightarrow A_k > 0$ and let (r_k, p_k) $(r_k = \rho(A_k), p_k > 0, ||p_k|| = 1)$ eigenpair
 - $\Rightarrow \{p_k\}_{k=1}^{\infty}$ is a bounded set as contained in the unit 1-sphere in \mathbb{R}^n
 - Convergence: each bounded sequence in \mathbb{R}^n has a convergent subsequence

$$\Rightarrow \{p_k\}_{k=1}^{\infty} \text{ has a convergent subsequence, } p_{k_i} > 0 \text{ and } \|p_{k_i}\| = 1 \\ \Rightarrow \{p_{k_i}\}_{k_i=1}^{\infty} \to z \text{ where } z \ge 0$$

• Take limitations: $r_k = \lim_{t \to \infty} \|A_k^t\|^{1/t} \Rightarrow 0 \le A < A_1, \ \rho(A) \le \rho(A_1)$ $\Rightarrow A_1 > A_2 > \dots > A \Rightarrow r_1 > r_2 > \dots > r \ (r = \rho(A)), \ \{r_k\}_{k=1}^{\infty} \text{ is a monotonic sequence of positive numbers bounded by } r$ $\Rightarrow \lim_{k \to \infty} r_k = r^* \text{ exists and } r^* \ge r, \lim_{k_i \to \infty} r_{k_i} = r^* \ge r$ $\Rightarrow \lim_{k \to \infty} A_k = A \text{ implies } \lim_{k_i \to \infty} A_{k_i} = A$ $\Rightarrow Az = \lim_{k_i \to \infty} A_{k_i} p_{k_i} = \lim_{i \to \infty} r_{k_i} p_{k_i} = r^* z \Rightarrow r^* \in \operatorname{spec}(A) \Rightarrow r^* \le r$ $\Rightarrow r^* = r \text{ and } Az = rz \text{ with } z \ge 0 \text{ and } z \neq 0$

• Proof of 3) and 4) for irreducible matrices

- $\rho(A)$ is simple: $r = \rho(A)$, let $B = (I + A)^{n-1} > 0$ and $\nu = \rho(B)$ $\Rightarrow \lambda \in \operatorname{spec}(A) \Leftrightarrow (1 + \lambda)^{n-1} \in \operatorname{spec}(B)$, $\operatorname{algmulti}_A(\lambda) = \operatorname{algmulti}_B((1 + \lambda)^{n-1})$ $\Rightarrow \nu = \max_{\lambda \in \operatorname{spec}(A)} |1 + \lambda|^{n-1} = \{\max_{\lambda \in \operatorname{spec}(A)} |1 + \lambda|\}^{n-1} = (1 + r)^{n-1}$ $\Rightarrow \operatorname{algmulti}_A(r) = 1 \Leftrightarrow \operatorname{algmulti}_B(\nu) = 1$. • Positive eigenvector: (r, x) is eigenpair of $A \Leftrightarrow (\nu, x)$ is eigenpair of B
 - $\Rightarrow x > 0$
 - $\Rightarrow r > 0$; otherwise Ax = 0 impossible $\Leftarrow A \ge 0$, $x > 0 \Rightarrow Ax > 0$

• Proof of 5)

- By definition of primitive matrix $B = A^k > 0 \Rightarrow \lambda \in \operatorname{spec}(A) \Leftrightarrow \lambda^k \in \operatorname{spec}(B)$
 - Suppose $|\lambda_1| = 1$ and $\lambda_1 \neq \rho(A) \Rightarrow \lambda_1^k \in \operatorname{spec}(B)$ $\Rightarrow |\lambda_1^k| = 1$ and $\operatorname{spec}(B)$ has two eigenvalues on the spectral circle contradict with the result for positive matrix only one eigenvalue $\rho(A)$ on the spectral circle
- $\bullet \ \Rightarrow \ {\rm eigenvalue} \ \rho(A) > |\mu|$ for all other eigenvalues μ





(a) The matrix A_3 is reducible: its dominant eigenvalue is 0 and so is its other eigenvalue.

(b) The matrix A₂ is irreducible but not primitive: its dominant eigenvalues +1 is not strictly larger, in magnitude, than the other eigenvalue -1.

 $\lambda = \overline{1}$



(c) The matrix A_4 is primitive: its dominant eigenvalue +1 is strictly larger than the other eigenvalue -1/2.

Figure 9: Example 2-dimensional non-negative matrices and their properties

28 / 32

Perron-Frobenius Theory

• Powers of non-negative matrices

Theorem

Let $A \in \mathbb{R}^{n \times n}$, $n \ge 2$ be non-negative with dominant eigenvalue λ and the right and left eigenvectors are denoted by v and w of λ , $v^{\top}w = 1$. If λ is simple and strictly larger in magnitude than all other eigenvalues, then we have

$$\lim_{k \to \infty} \frac{A^k}{\lambda^k} = v w^\top \tag{18}$$

Proof.

$$\begin{split} \lambda \text{ is simple and strictly larger } &\lambda = T \begin{bmatrix} \lambda & 0_{1 \times n-1} \\ 0_{n-1 \times 1} & B \end{bmatrix} T^{-1} \text{ and } \rho(B/\lambda) < 1 \\ \Rightarrow &\lim_{k \to \infty} B^k / \lambda^k = 0 \Rightarrow \lim_{k \to \infty} (\frac{A}{\lambda})^k = T (\lim_{k \to \infty} \begin{bmatrix} 1^k & 0_{1 \times n-1} \\ 0_{n-1 \times 1} & B^k \end{bmatrix}) T^{-1} = \\ T (\lim_{k \to \infty} \begin{bmatrix} 1 & 0_{1 \times n-1} \\ 0_{n-1 \times 1} & 0_{n-1 \times n-1} \end{bmatrix}) T^{-1} = v w^{\top}, v \text{ is the first column of } T \text{ and } w \text{ is the first row of } T^{-1}. \end{split}$$

- Row-stochastic matrices (Let A be a row-stochastic matrix and let G be its associated digraph)
 - ${\, \bullet \,}$ the eigenvalue 1 is simple and all other eigenvalues $|\mu|<1$
 - $\lim_{k \to \infty} A^k = 1_n w^\top$ for w > 0 and $1_n^\top w = 1$
 - ${\ensuremath{\, \bullet \,}} G$ is an aperiodic strongly-connected graph

- Row-stochastic matrices (Let A be a row-stochastic matrix and let G be its associated digraph)
 - ${\, \bullet \,}$ the eigenvalue 1 is simple and all other eigenvalues $|\mu|<1$
 - $\lim_{k \to \infty} A^k = 1_n w^\top$ for w > 0 and $1_n^\top w = 1$
 - ${\ensuremath{\, \bullet \,}} G$ is an aperiodic strongly-connected graph
- If the previous conditions are satisfied, then
 - the solution of x(k+1) = Ax(k) satisfies $\lim_{k \Rightarrow \infty} x(k) = w^{\top}x(0)1_n$
 - if additionally A is doubly-stochastic, then $w = \frac{1}{n} \mathbf{1}_n$ so that

$$\lim_{k \to \infty} x(k) = \frac{\mathbf{1}_n^\top x(0)}{n} \mathbf{1}_n = \operatorname{average}(x(0)) \mathbf{1}_n$$
(19)

• Discrete-time averaging systems

- Background and application examples
- Analysis intuition for convergence (consensus)
- Conditions to ensure consensus and average consensus

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Thank You! Q&A

C. Zhao (ZJU)

Discrete-time Averaging Systems

August 2, 2021 32 / 32

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